

A New Mathematical Programming Framework for Facility Layout

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The Facility Layout (or Floorplanning) Problem

- Find the optimal positions for a given set of N rectangular departments of *fixed area* within a rectangular facility of *fixed area*.
- All the dimensions may be given or left undetermined.
- The objective is to minimize (according to some norm, e.g. l_1, l_2) the distances between pairs of departments that have a nonzero connection “cost”.

Applications ?

- Hospital layout
- Service center layout
- VLSI placement and design
- etc.

But : Like many optimization problems from practical applications, the facility layout problem is “hard” (NP-hard).

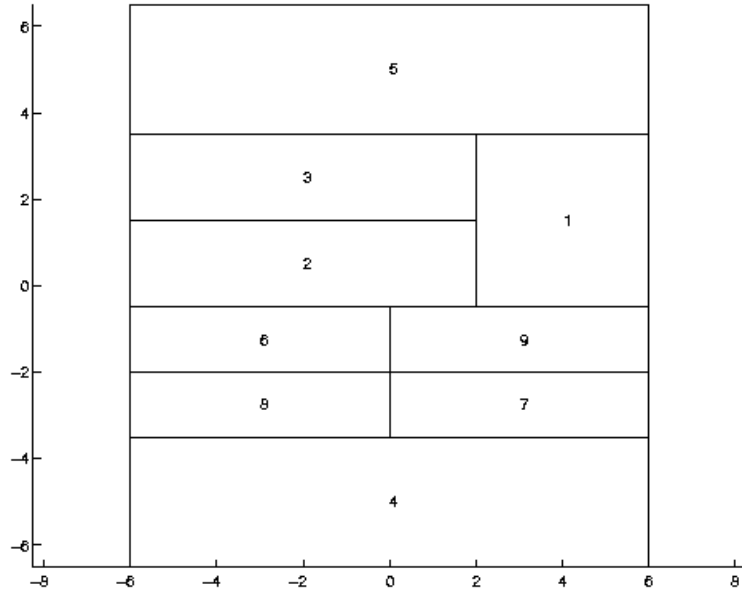
Small Example

| Dept | Area |
|------|------|
| 1 | 16 |
| 2 | 16 |
| 3 | 16 |
| 4 | 36 |
| 5 | 36 |
| 6 | 9 |
| 7 | 9 |
| 8 | 9 |
| 9 | 9 |

| i | j | C_{ij} | i | j | C_{ij} | i | j | C_{ij} |
|-----|-----|----------|-----|-----|----------|-----|-----|----------|
| 1 | 4 | 5 | 2 | 9 | 1 | 4 | 7 | 4 |
| 1 | 5 | 5 | 3 | 4 | 2 | 5 | 6 | 3 |
| 1 | 9 | 1 | 3 | 5 | 2 | 5 | 9 | 4 |
| 2 | 4 | 3 | 3 | 9 | 1 | 6 | 9 | 2 |
| 2 | 5 | 3 | 4 | 6 | 4 | 7 | 9 | 1 |

Facility is fixed to be a 12 x 13 rectangle

Possible Layout for Example



Total Euclidean cost: 229.71

Motivation for this Work

- Exact mixed integer programming approaches only work for problems with less than 10 departments.
- Most other approaches in the literature are based on heuristic search methods.
- ➔ We present a new two-stage framework based on mathematical programming models, and inspired by a *convex global relaxation* of the layout problem.

Outline of the Proposed New Framework

- ✓ The first model is a *convex relaxation* of the layout problem (to find a good starting point);
 - ✓ The second model is an *exact formulation* of the problem as a mathematical program with equilibrium constraints (MPEC).
- ➔ Both models can be *solved efficiently* using widely available non-linear optimization software.

The NLT method

(van Camp, Carter & Vannelli, 1991)

The (non-convex) vCCV model

$$\begin{aligned}
 & \min_{(x_i, y_i), h_i, w_i, h_F, w_F} \sum_{1 \leq i < j \leq N} c_{ij} \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2} \\
 \text{s. t. } & \begin{cases} |x_i - x_j| - \frac{1}{2}(w_i + w_j) \geq 0 & \text{if } |y_i - y_j| - \frac{1}{2}(h_i + h_j) < 0, \forall i \neq j, \\ |y_i - y_j| - \frac{1}{2}(h_i + h_j) \geq 0 & \text{if } |x_i - x_j| - \frac{1}{2}(w_i + w_j) < 0, \forall i \neq j, \end{cases} \\
 & \begin{cases} \frac{1}{2} w_F - (x_i + \frac{1}{2} w_i) \geq 0 & \forall i, & \frac{1}{2} h_F - (y_i + \frac{1}{2} h_i) \geq 0 & \forall i, \\ (x_i - \frac{1}{2} w_i) + \frac{1}{2} w_F \geq 0 & \forall i, & (y_i - \frac{1}{2} h_i) + \frac{1}{2} h_F \geq 0 & \forall i, \end{cases} \\
 & \begin{cases} \min(w_i, h_i) - l_i^{\min} \geq 0 & \forall i, & l_i^{\max} - \min(w_i, h_i) \geq 0 & \forall i, \\ \min(w_F, h_F) - l_F^{\min} \geq 0, & & l_F^{\max} - \min(w_F, h_F) \geq 0, & \end{cases} \\
 & h_i w_i = a_i \quad \forall i
 \end{aligned}$$

where l_i^{\min} , l_i^{\max} , l_F^{\min} and l_F^{\max} are given bounds on the dimensions.

NLT: Three-Stage Approach

- (1) Evenly distribute the centres of the departments inside the facility;
- (2) Reduce the overlap between departments;
- (3) Determine the final solution by solving the vCCV model.

Stages 1 and 2 are (non-convex) relaxations of the vCCV model which approximate the departments by circles.

The Stage-2 model of NLT

$$\min_{(x_i, y_i), h_F, w_F} \sum_{1 \leq i < j \leq N} c_{ij} d_{ij}$$

$$\text{subject to } d_{ij} \geq r_i + r_j \quad \forall i \neq j$$

$$\frac{1}{2} w_F \geq x_i + r_i, \quad \frac{1}{2} w_F \geq r_i - x_i \quad \forall i$$

$$\frac{1}{2} h_F \geq y_i + r_i, \quad \frac{1}{2} h_F \geq r_i - y_i \quad \forall i$$

$$l_F^{\max} \geq \min(w_F, h_F) \geq l_F^{\min},$$

$$\text{where } d_{ij} := \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$$

and r_i is the radius of the circle for department i .

**ModCoAR:
The First Model of
the New Framework**

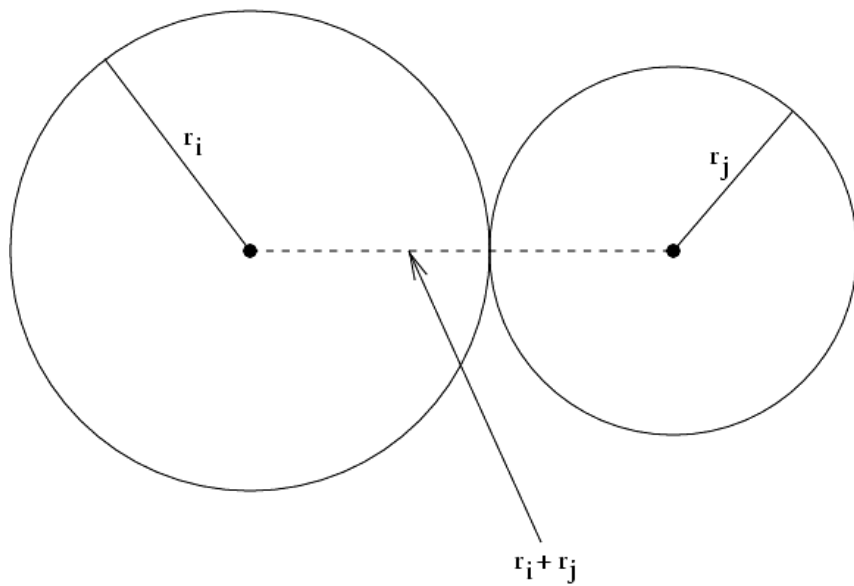
Steps to derive the ModCoAR model

- (1) Define the *target distance* concept for circles with varying radii;
- (2) Enforce the target distances using a *repeller* term in the objective function;
- (3) Analyse and *convexify* the result
 - ➔ Concept of *generalized target distances*;
- (4) Add a *barrier term* (for ease of computation).

For convenience, we work with the squares of the distances:

$$D_{ij} := d_{ij}^2 = (x_i - x_j)^2 + (y_i - y_j)^2$$

Target distances concept



Hence the target distance between circles i and j is

$$t_{ij} := \alpha \cdot (r_i + r_j)^2 \quad \text{for some } \alpha > 0.$$

Attractor-Repeller Paradigm

For each pair i,j of modules, the distance minimizing term is viewed as an attractor:

$$c_{ij} \cdot D_{ij}, \quad c_{ij} \geq 0, \quad D_{ij} \geq 0$$

is minimized when $D_{ij} = 0$.

Enforcing the target distances

To counter this “attraction”, we enforce the target distances with *repeller* terms in the objective function:

$$f(z) := \frac{1}{z} - 1, \quad z > 0$$

and $z = \frac{D_{ij}}{t_{ij}}$, where t_{ij} is the target distance.

The (non-convex) AR model

$$\min_{(x_i, y_i), h_F, w_F} \sum_{1 \leq i < j \leq N} c_{ij} D_{ij} + \sum_{1 \leq i < j \leq N} f\left(\frac{D_{ij}}{t_{ij}}\right)$$

subject to ~~$d_{ij} \geq r_i + r_j \quad \forall i \neq j$~~

$$\frac{1}{2} w_F \geq x_i + r_i \quad \forall i$$

$$\frac{1}{2} h_F \geq y_i + r_i \quad \forall i$$

$$\frac{1}{2} w_F \geq r_i - x_i \quad \forall i$$

$$\frac{1}{2} h_F \geq r_i - y_i \quad \forall i$$

~~$$l_F^{\max} \geq \min(w_F, h_F) \geq l_F^{\min}$$~~

$$w_F^{\max} \geq w_F \geq w_F^{\min}$$

$$h_F^{\max} \geq h_F \geq h_F^{\min}$$

Examine the objective function

Rewrite the objective function:

$$\sum_{1 \leq i < j \leq N} c_{ij} D_{ij} + \sum_{1 \leq i < j \leq N} \left(\frac{t_{ij}}{D_{ij}} - 1 \right) = \sum_{1 \leq i < j \leq N} \left(c_{ij} D_{ij} + \frac{t_{ij}}{D_{ij}} - 1 \right)$$

and since the sum of convex functions is convex,
we ask:

When is the term $c_{ij} D_{ij} + \frac{t_{ij}}{D_{ij}} - 1$ convex?

Fact: Let

$$g : \mathfrak{R}^4 \rightarrow \mathfrak{R}, \quad g(x_1, x_2, y_1, y_2) = c z + \frac{t}{z} - 1,$$

where

$$c > 0, t > 0 \text{ and } z > 0, z = (x_1 - x_2)^2 + (y_1 - y_2)^2.$$

Then the following statements hold for g :

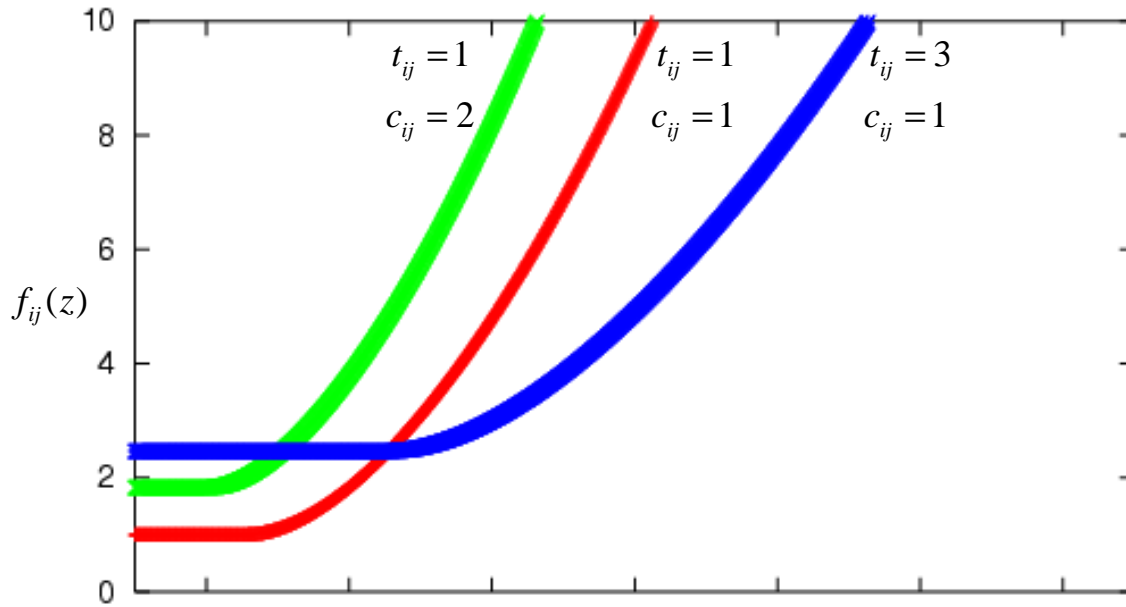
1. If $z \geq \sqrt{\frac{t}{c}}$ then the Hessian of g is positive semidefinite.
2. If $z = \sqrt{\frac{t}{c}}$ then the gradient of g is zero.

Define a new (convex!) function

For $c_{ij} > 0$, $t_{ij} > 0$, and $D_{ij} = (x_i - x_j)^2 + (y_i - y_j)^2$, we define the convex, continuously differentiable piecewise function

$$f_{ij}(x_i, x_j, y_i, y_j) := \begin{cases} c_{ij} D_{ij} + \frac{t_{ij}}{D_{ij}} - 1, & D_{ij} \geq \sqrt{\frac{t_{ij}}{c_{ij}}} \\ 2 \sqrt{c_{ij} t_{ij}} - 1, & 0 \leq D_{ij} < \sqrt{\frac{t_{ij}}{c_{ij}}} \end{cases}$$

Graph of f_{ij}



The *convex* CoAR model

$$\min_{(x_i, y_i), h_F, w_F} \sum_{1 \leq i < j \leq N} f_{ij}(x_i, x_j, y_i, y_j)$$

$$\text{subject to } \frac{1}{2} w_F \geq x_i + r_i \quad \forall i \in M$$

$$\frac{1}{2} h_F \geq y_i + r_i \quad \forall i \in M$$

$$\frac{1}{2} w_F \geq r_i - x_i \quad \forall i \in M$$

$$\frac{1}{2} h_F \geq r_i - y_i \quad \forall i \in M$$

$$w_F^{\max} \geq w_F \geq w_F^{\min}$$

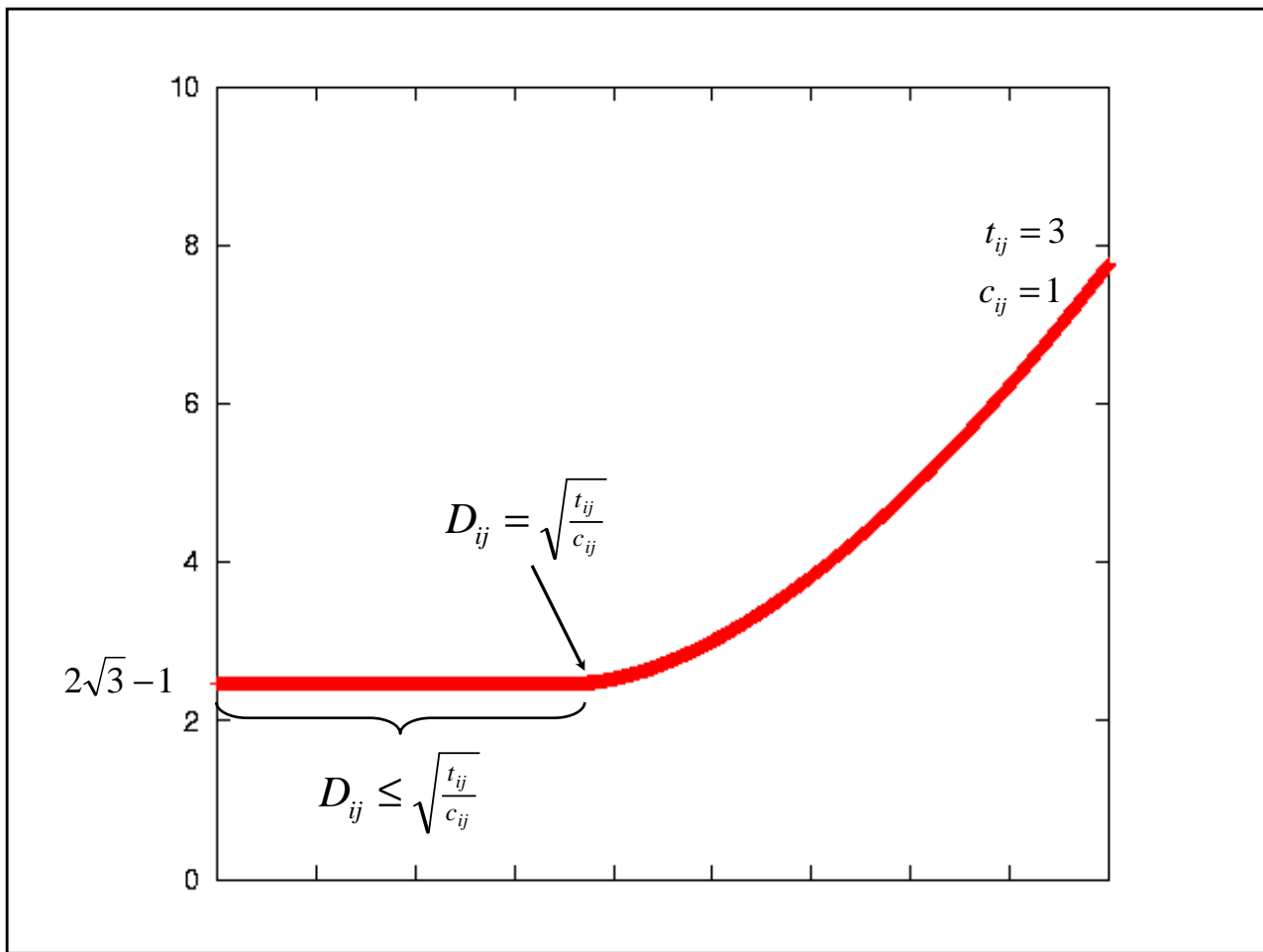
$$h_F^{\max} \geq h_F \geq h_F^{\min}$$

When is the value of f_{ij} minimum?

- We deduce from the structure of f_{ij} that its minimum value is attained for all positions of the departments i and j for which

$$D_{ij} \leq \sqrt{\frac{t_{ij}}{c_{ij}}}.$$

- This includes $D_{ij}=0$ (complete overlap!).



Since we want to minimize overlap, what we really want is a layout for which

$$D_{ij} \approx \sqrt{\frac{t_{ij}}{c_{ij}}}.$$

For such a point, we have

D_{ij} proportional to t_{ij} ;

hence our original target distances are still enforced.

Generalized Target Distances

If we define

$$T_{ij} := \sqrt{\frac{t_{ij}}{c_{ij} + \varepsilon}}, \quad \varepsilon > 0 \quad \text{"small"}$$

then we can think of T_{ij} as a

generalized target distance

for the departments i and j .

This “new” target distance takes both t_{ij} and c_{ij} into account.

Practical Interpretation of T_{ij}

$$T_{ij} := \sqrt{\frac{t_{ij}}{c_{ij} + \varepsilon}}$$

- If c_{ij} is small, then departments i and j are likely to be placed far apart in the layout, so the corresponding T_{ij} can be large;
- If c_{ij} is large, then the opposite reasoning applies, and T_{ij} can be small;
- But T_{ij} also takes t_{ij} into account!

How to achieve T_{ij} ?

Add to the objective function a term of the form

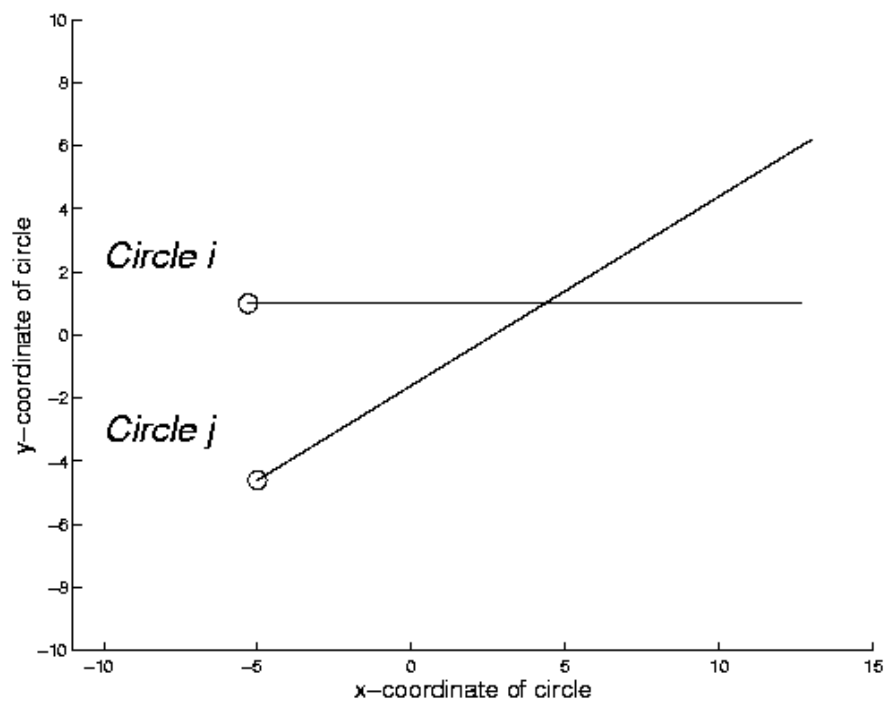
$$-K \ln\left(\frac{D_{ij}}{T_{ij}}\right)$$

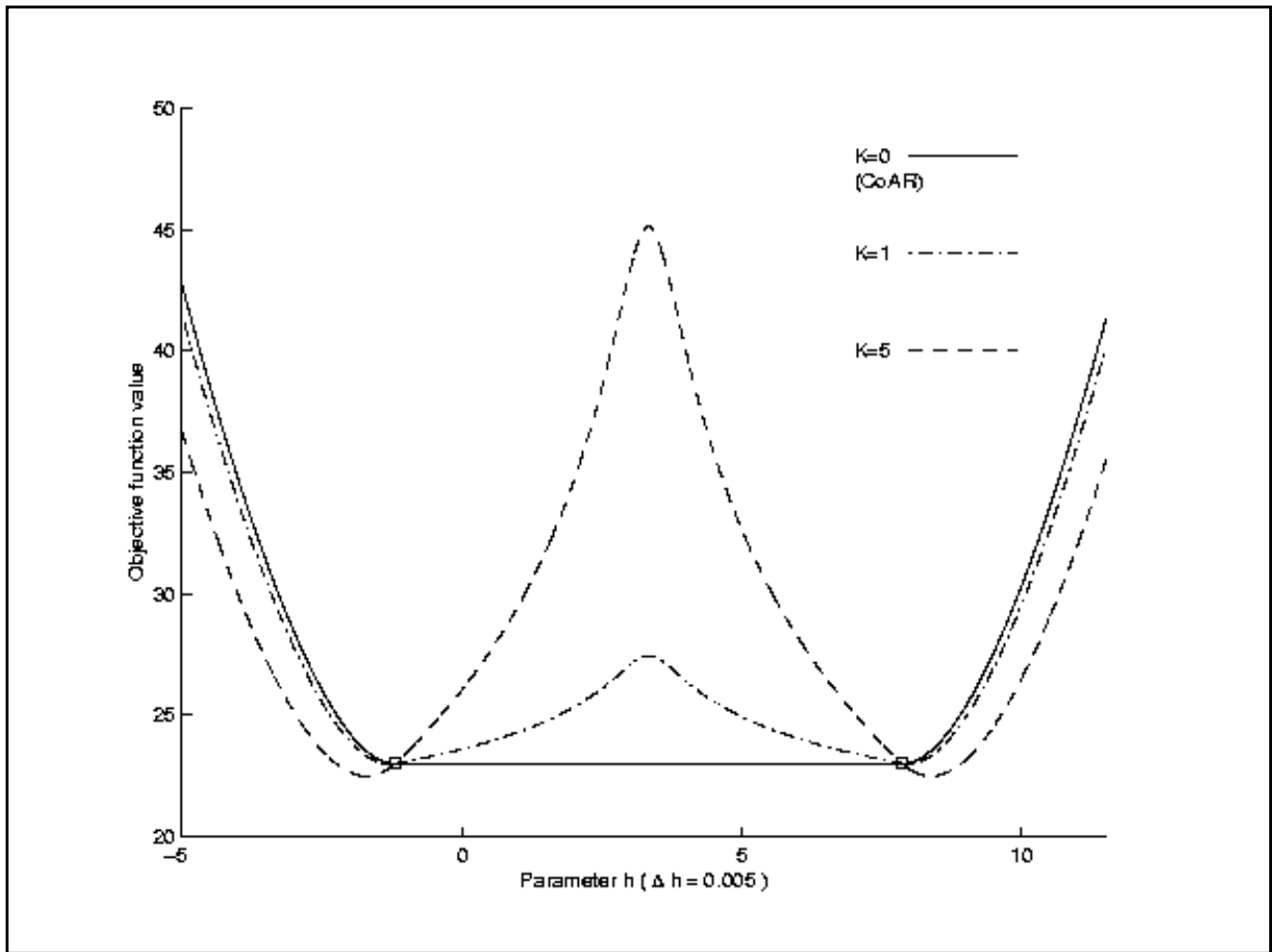
for each pair of departments.

The resulting function has minima that satisfy

$$D_{ij} \approx T_{ij}$$

A Glimpse of the 4D Function...





The ModCoAR model

$$\min_{(x_i, y_i), h_F, w_F} \sum_{1 \leq i < j \leq N} F_{ij}(x_i, x_j, y_i, y_j) - K \ln(D_{ij} / T_{ij})$$

$$\text{subject to} \quad \frac{1}{2} w_F \geq x_i + r_i, \quad \frac{1}{2} h_F \geq y_i + r_i \quad \forall i$$

$$\frac{1}{2} w_F \geq r_i - x_i, \quad \frac{1}{2} h_F \geq r_i - y_i \quad \forall i$$

$$w_F^{\max} \geq w_F \geq w_F^{\min}, \quad h_F^{\max} \geq h_F \geq h_F^{\min}$$

where

$$F_{ij}(x_i, x_j, y_i, y_j) := \begin{cases} c_{ij} D_{ij} + \frac{t_{ij}}{D_{ij}} - 1, & D_{ij} \geq T_{ij} \\ 2 \sqrt{c_{ij} t_{ij}} - 1, & 0 \leq D_{ij} < T_{ij} \end{cases}$$

BPL :
The Second Model of
the New Framework

Non-overlap constraints

$$\left| x_i - x_j \right| - \frac{1}{2}(w_i + w_j) \geq 0 \quad \text{if} \quad \left| y_i - y_j \right| - \frac{1}{2}(h_i + h_j) < 0$$

$$\left| y_i - y_j \right| - \frac{1}{2}(h_i + h_j) \geq 0 \quad \text{if} \quad \left| x_i - x_j \right| - \frac{1}{2}(w_i + w_j) < 0$$

Note that these constraints are *disjunctive*...

... and therefore can be written as

$$\frac{1}{2}(w_i + w_j) - |x_i - x_j| \leq 0 \quad \text{or} \quad \frac{1}{2}(h_i + h_j) - |y_i - y_j| \leq 0$$

which is equivalent to

$$\min \left\{ \frac{1}{2}(w_i + w_j) - |x_i - x_j|, \frac{1}{2}(h_i + h_j) - |y_i - y_j| \right\} \leq 0$$

New variables

For each pair of departments, introduce two new variables

$$X_{ij}, Y_{ij}$$

and let

$$X_{ij} \geq \frac{1}{2}(w_i + w_j) - |x_i - x_j|, \quad X_{ij} \geq 0$$

$$Y_{ij} \geq \frac{1}{2}(h_i + h_j) - |y_i - y_j|, \quad Y_{ij} \geq 0$$

Equilibrium Constraints

Then

$$\min \left\{ \frac{1}{2} (w_i + w_j) - |x_i - x_j|, \frac{1}{2} (h_i + h_j) - |y_i - y_j| \right\} \leq 0$$

is equivalent to

$$X_{ij} Y_{ij} = 0$$

MPEC Formulation

(Math. Prog. With Equilibrium Constraints)

$$\min_{(x_i, y_i), h_i, w_i, h_F, w_F} \sum_{1 \leq i < j \leq N} c_{ij} \delta_{ij}$$

$$\text{s. t. } X_{ij} \geq \frac{1}{2}(w_i + w_j) - |x_i - x_j| \quad , \quad Y_{ij} \geq \frac{1}{2}(h_i + h_j) - |y_i - y_j|$$

$$X_{ij} Y_{ij} = 0 \quad , \quad X_{ij} \geq 0 \quad , \quad Y_{ij} \geq 0 \quad , \quad \forall 1 \leq i < j \leq N$$

$$h_i w_i = a_i \quad \forall i \quad (\text{area constraints})$$

plus: “fit-in-the-facility” constraints and bound constraints
(all linear)

and $\delta_{ij}(x_i, x_j, y_i, y_j)$ is the desired norm (l_1, l_2, \dots) .

Computational Results

Solution methodology

- We solve both models using the software package MINOS.
- For ModCoAR, because of the linearity of the constraints, convergence is generally *superlinear*.

Solution methodology (ctd)

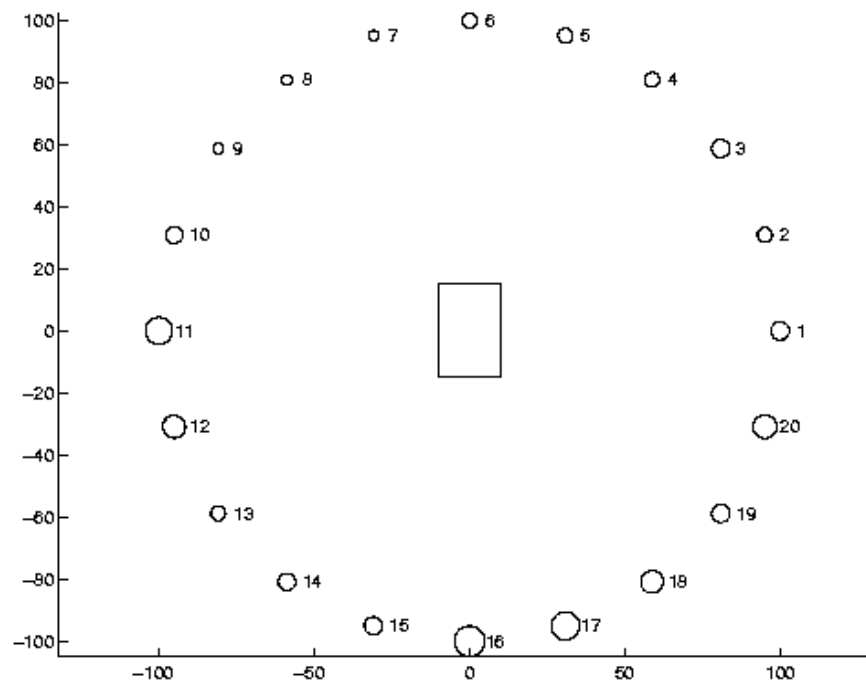
- We chose to set

$$K = \sum_{1 \leq i < j \leq N} c_{ij}$$

so that K clearly dominates the c_{ij} 's.

- MINOS requires an initial configuration to start the (iterative) algorithm for solving ModCoAR.

Choice of initial configuration



Solving the MPEC using MINOS

- The complementarity constraints

$$X_{ij} Y_{ij} = 0 \quad , \quad X_{ij} \geq 0 \quad , \quad Y_{ij} \geq 0$$

imply that **no** strictly feasible point exists.

This causes MINOS to fail...

- Thus we apply a penalty-type approach to the above constraints.

BPL Model

$$\min_{(x_i, y_i), h_i, w_i, h_F, w_F} \sum_{1 \leq i < j \leq N} c_{ij} \delta_{ij} + K \cdot X_{ij} Y_{ij}$$

$$\text{s. t. } X_{ij} \geq \frac{1}{2}(w_i + w_j) - |x_i - x_j|, \quad Y_{ij} \geq \frac{1}{2}(h_i + h_j) - |y_i - y_j|$$

$$X_{ij} \geq 0, \quad Y_{ij} \geq 0, \quad \forall 1 \leq i < j \leq N$$

$$h_i w_i = a_i \quad \forall i$$

plus: “fit-in-the-facility” constraints and bound constraints
(all linear)

Aspect ratio constraints

- The aspect-ratio for department i is defined as

$$\beta_i := \frac{\max\{h_i, w_i\}}{\min\{h_i, w_i\}}$$

- Bounding (above) the aspect ratio ensures that no departments are excessively narrow in the layout.

Aspect ratio constraints (ctd)

- We enforce the bound β_i^* on β_i by adding to BPL the constraints

$$\beta_i w_i \geq h_i, \quad \beta_i h_i \geq w_i, \quad \beta_i^* \geq \beta_i.$$

Classical example :

Armour & Buffa problem (1963)

- Large problem (20 departments) – beyond all previous mathematical programming approaches (mixed integer programming).
- Each run of our algorithm requires approximately 18 seconds of CPU time (300 MHz SunSPARC).
- We can compare our framework using the rectilinear norm with the most recent results in the literature (Tate & Smith'95 -- genetic algorithm).

Experiments with the Armour & Buffa problem (1)

First we set no aspect ratio constraints, only a lower bound of 2 on all heights and widths.

We found a layout with cost 4230.6
and aspect ratio 6.67

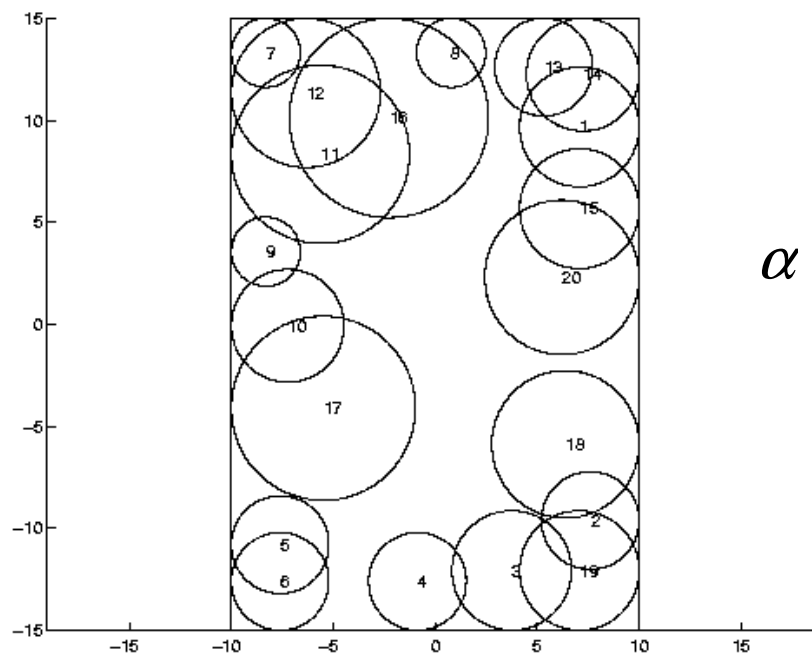
In TS'95, the best layout with aspect ratio bounded by 7 has cost 5255.0

Experiments with the Armour & Buffa problem (2)

Then we started setting aspect ratio constraints:

| β_i^* | TS'95 | New framework |
|-------------|--------|------------------|
| 5 | 5524.7 | 4591.3 |
| 4 | 5743.1 | 4786.4 |
| 3 | 5832.6 | 5140.1 |
| 2 | 6171.1 | 5224.7 |

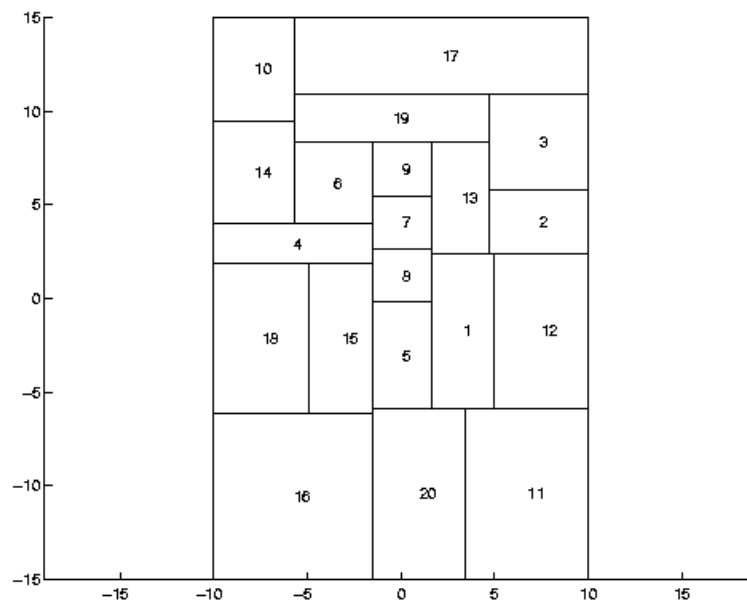
Best layout with $\beta_i \leq 4$



$$\alpha = 1.532$$

Optimal solution of ModCOAR

Best layout with $\beta_i \leq 4$ (ctd)



Total cost 4786.4 (versus 5743.1 in TS'95)

On-going and Future Research

- Apply this framework to the MCNC macro-cell layout problems, and compare the results with other methods.
- Investigate more thoroughly the role of α in the model.
- Improve the solution methodology; in particular, apply a nonlinear programming solver that directly tackles the MPEC formulation in spite of the lack of a strictly feasible point.

References

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