

# Isostatic Structures: Using Richard Rado's Matroid Matchings

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# Outline

- 1 Main Points
- 2 Basics and Context
- 3 Semi-simplicial Maps
- 4 Shelling
- 5 Freely Shellable Maps
- 6 Partitions of the Vertex Set
- 7 Finale

# Dedication

I would like to dedicate this talk to two persons,  
both of whom are architects and engineers.

# Dedication

To *Janos Baracs*,



# Dedication

To *Janos Baracs*,



instigator and cofounder of  
the research group *Topologie Structurale*,  
who learned projective geometry  
from his high school math teacher in Budapest,  
and who introduced Ivo Rosenberg and myself to  
three dimensional space and rigidity  
during a workshop for members of the  
Centre de recherches mathématiques  
in January 1973, over 38 years ago,

# Dedication

... posing, among other problems:

- to characterize generically 3-isostatic graphs
- to predict special positions of non-rigidity for generically 3-isostatic graphs,
- to specify the correct placements of cross-braces in grid frameworks.
- to analyze the rigidity of tensegrity frameworks.
- to analyze the relation between stresses and lifting of plane polyedral frameworks.
- to develop a theory of periodic filling of space by copies of one or more associated zonohedra.

# Dedication

To *Richard Gage*,



# Dedication

To *Richard Gage*,



founder and leading member of  
the association *Architects and Engineers for 911 Truth*,  
who has brought a new level of intelligent and systematic inquiry,  
a new level of organization and energetic public engagement,  
to the quest for an independent inquiry into  
the state crimes of 11/9/2001  
and into this decade of their rain of miserable consequences.





# Dedication

Everything you ever wanted to know  
about the 9/11 conspiracy theory  
in under 5 minutes.

<http://www.informationclearinghouse.info/article29110.htm>

*(surely the central rigidity problem of our era)*

# Dedication

With special thanks to Walter Whiteley and Bob Connelly,  
Ileana Streinu and Tibor Jordán,  
who have so energetically  
kept this beautiful subject alive and well,  
expanding its horizons,  
training the researchers of this new generation,  
and making it possible for us to be together today.

# Main Points

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- (3) We offer a strengthened conjecture:  
**Conjecture:** A graph is generically  $d$ -isostatic if and only if it has a *freely-shellable* semi-simplicial map to the  $d$ -simplex.

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- (4) We investigate further restrictions of the class of maps to maps that are fewer in number and easier to construct: maps whose vertex packets are *broken paths*.

# Generically Isostatic Graphs

A graph  $G(V, E)$  is generically  $d$ -isostatic if and only if it is edge-minimal among graphs that are rigid in some (and therefore in almost every) position in real Euclidean or projective space of dimension  $d$ .



# Generically Isostatic Graphs

We shall deal only with **generic behavior** of graphs as structures, so we will speak simply of “ $d$ -isostatic” graphs, dropping the adjective “generic”.

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# $d$ -Isostatic Graphs

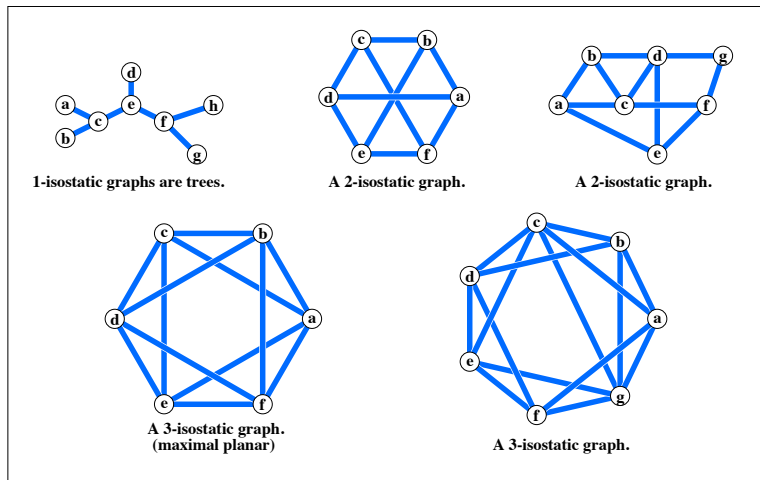


Figure:  $d$ -Isostatic graphs, for  $d = 1, 2, 3$ .

## Definition

A *semi-simplicial map*  $f : G(V, E) \rightarrow K_{d+1}$ ,

where

$$K_{d+1} = K(I, J)$$

and

$$I = \{1, 2, \dots\}, \quad J = \{12, 13, \dots\},$$

consists of a pair of maps

$$f_0 : V \rightarrow I, f_1 : E \rightarrow J,$$

that preserve incidence.



# Definition

That is, an edge  $e = ab$  whose vertices  $a$  and  $b$  have **distinct values**  $f_0(a) = i$ ,  $f_0(b) = j$  in  $I$  must be sent by  $f_1$  to  $ij \in J$ .

We call such an edge  $e = ab$  an  **$ij$ -bridge**.

# Definition

An edge  $e = ab$  whose end vertices go to **the same vertex**, say  $f_0(a) = i = f_0(b)$ , must be sent to an edge  $ij$  of  $K$  incident to  $i$ .

We call such an edge  $e = ab$  a *loop at  $i$  toward  $j$* .

# Definition

The subset  $f_0^{-1}(i)$ , for any vertex  $i \in I$ ,  
we call the  $i^{\text{th}}$  *vertex packet* of  $f$ , denoted  $V_i$ .

# Definition

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$(\mathcal{P}_0)$  *Edge independence*: The inverse image  $f_1^{-1}(ij)$ , denoted  $T_{ij}$ , of any edge  $ij$  of  $K$  is a tree spanning the union  $V_i \cup V_j$  of its two related vertex packets.

# Definition

(the combined statement:)

A *semi-simplicial map*  $f : G(V, E) \rightarrow K_{d+1}(I, J)$ , consists of a pair of maps  $f_0 : V \rightarrow I$ ,  $f_1 : E \rightarrow J$ , that preserve incidence, and ...

( $\mathcal{P}_0$ ) *Edge independence*: The inverse image  $f_1^{-1}(ij)$ , denoted  $T_{ij}$ , of any edge  $ij$  of  $K$  is a tree spanning the union  $V_i \cup V_j$  of its two related vertex packets.

# Visual Representation of Maps

Semi-simplicial maps have very satisfactory visual representations, using *colors* taken from a standard edge-coloring of  $K_{d+1}$  to specify the images of each edge.

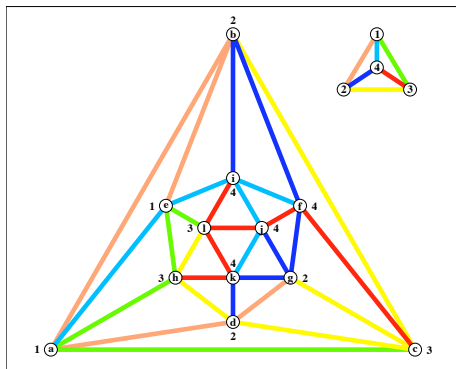


Figure: A  $d$ -isostatic graph, with semi-simplicial map.

# Visual Representation of Maps

The tree-decomposition is then easily comprehended.

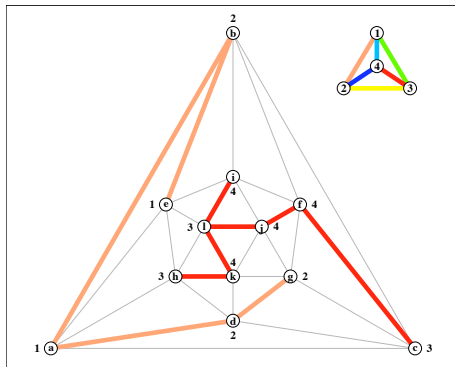


Figure: The trees  $T_{12}$  and  $T_{34}$ .



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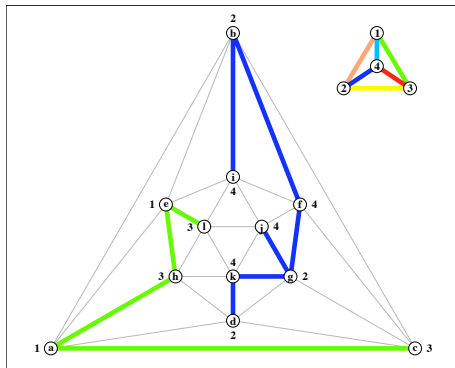


Figure: The trees  $T_{13}$  and  $T_{24}$ .

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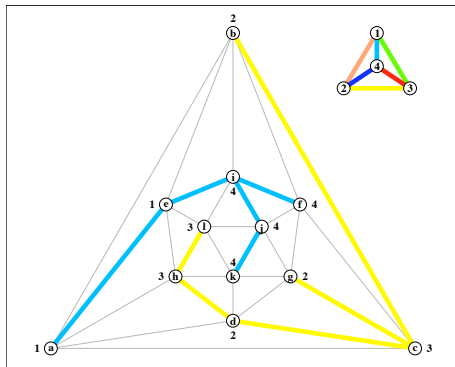


Figure: The trees  $T_{14}$  and  $T_{24}$ .

# Visual Representation of Maps

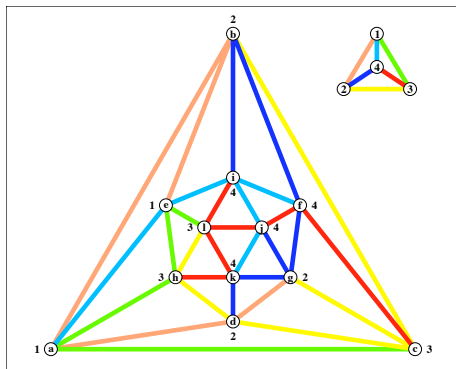


Figure: All together now!.

# Path Connectivity

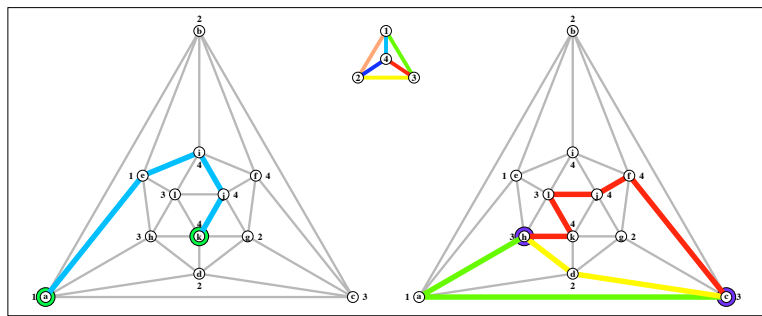


Figure: Paths between vertices having distinct/identical images.

If  $a$  and  $b$  have distinct images  $i, j$  under  $f_0$ ,  
then  $a$  and  $b$  are connected along a unique path in the tree  $T_{ij}$ .

# Path Connectivity

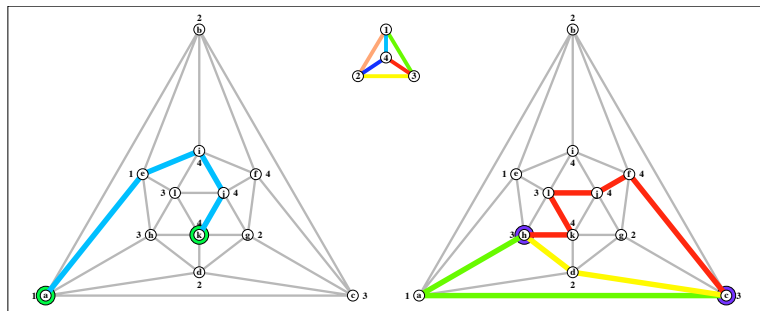


Figure: Paths between vertices having distinct/identical images.

If  $a$  and  $b$  have the same image  $i$  under  $f_0$ , then they are connected along unique paths in each of the  $d$  trees  $T_{ij}$ , for  $j \neq i$ .

# Shelling

A vertex packet can be **shelled** if there is a sequence of monochromatic cuts that reduces it to a subgraph with no edges.

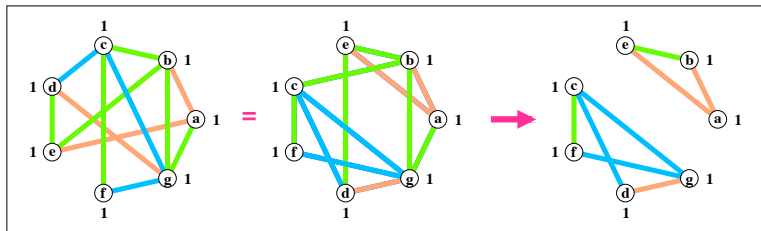


Figure: A sequence of monochromatic cuts.

# Special Placement

In the special position given by a semi-simplicial map, any external equilibrium test load applied at two vertices  $a, b$  is uniquely resolvable.

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In the special position given by a semi-simplicial map, any external equilibrium test load applied at two vertices  $a, b$  is uniquely resolvable.

If  $f(a) \neq f(b)$ , the external load is resolved (and uniquely so) along the path between  $a$  and  $b$  in the tree  $T_{ij}$ , all those edges being *collinear* along the line  $i \vee j$ .



# Special Placement

In the special position given by a semi-simplicial map, any external equilibrium test load applied at two vertices  $a, b$  is uniquely resolvable.

If  $f(a) = f(b) = i$ , the external load can be uniquely represented as a sum of  $d + 1$  equilibrium loads applied to  $a, b$ , one in each of the (independent) directions  $i \vee j$  at  $i$ .

These individual loads are then uniquely resolvable along the paths from  $a$  to  $b$  in the trees  $T_{ij}$

# Theorem

*A graph  $G$  is generically  $d$ -isostatic graph if it has a shellable semi-simplicial map to the  $d$ -simplex.*

# Maps on Dependent Graphs

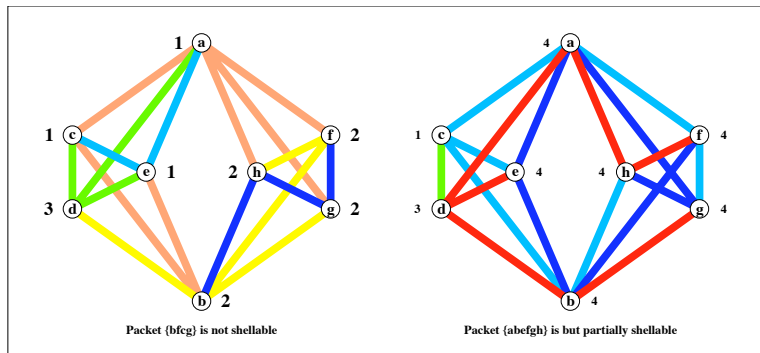


Figure: Non-shellable maps on a 3-dependent graph.

## Converse, $d = 2$

For  $d = 2$ : Any non-shellable map has an **obstacle to shelling** in the form of a set of 3 or more vertices co-spanned by sub-trees of two trees.

This is a **dependent** subgraph.

## Converse, $d = 2$

Theorem:

*A graph  $G$  is generically 2-isostatic graph  
if and only if  
it has a shellable semi-simplicial map to the triangle,  
if and only if  
all semi-simplicial maps to the triangle are shellable.*

Converse,  $d = 3$ ?

This is far from being the case in dimension 3.

A 3-isostatic graph may have many non-shellable maps to the tetrahedron.

## Converse, $d = 3$ ?

Existence of a non-shellable map establishes only that there is a subset  $Q$  of some vertex packet  $i$  that is spanned by sub-trees of any *pair* of the three trees  $T_{ij}$  for  $j \neq i$ .

Converse,  $d = 3$ ?

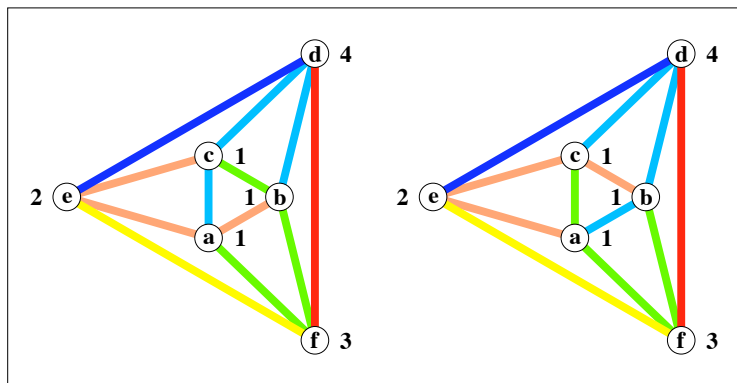


Figure: The packet  $V_1$  contains an obstacle to shelling.

These are the only two edge maps with this vertex map.



Converse,  $d = 3$ ?

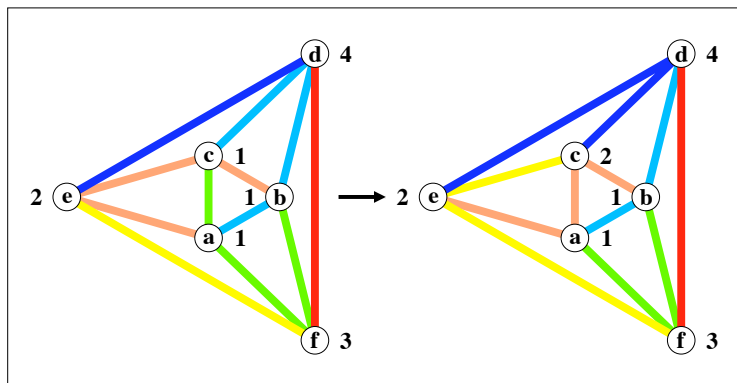


Figure: A change of one vertex image produces a shellable map.

This vertex map has a unique compatible edge map.

# Eliminate Obstacles - Eliminate Shelling

Perhaps the best way to deal with obstacles to shelling will be to look for maps in which obstacles cannot occur,

# Eliminate Obstacles - Eliminate Shelling

that is, those for which the vertex packets  
induce *independent* subgraphs,  
that is, cycle-free subgraphs, or forests.

# Eliminate Obstacles - Eliminate Shelling

These maps are *freely shellable*:

# Eliminate Obstacles - Eliminate Shelling

Simply proceed edge by edge,  
each single edge being a monochromatic cut!

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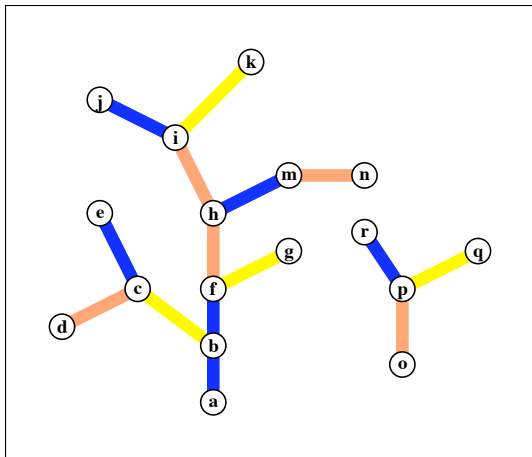


Figure: A forest as induced subgraph of packet  $V_2$ .

# Eliminate Obstacles - Eliminate Shelling

## Conjecture:

A graph  $G$  is generically 3-isostatic if and only if it has a semi-simplicial map to the tetrahedron in which all vertex packets induce subgraphs that are independent (ie: forests) as subgraphs of  $G$ .

# Eliminate Obstacles - Eliminate Shelling

There are four interesting classes of such maps:  
those in which the vertex packets induce:

$\mathcal{F}$  forests

$\mathcal{T}$  trees

$\mathcal{B}$  broken paths

$\mathcal{P}$  paths



# Eliminate Obstacles - Eliminate Shelling

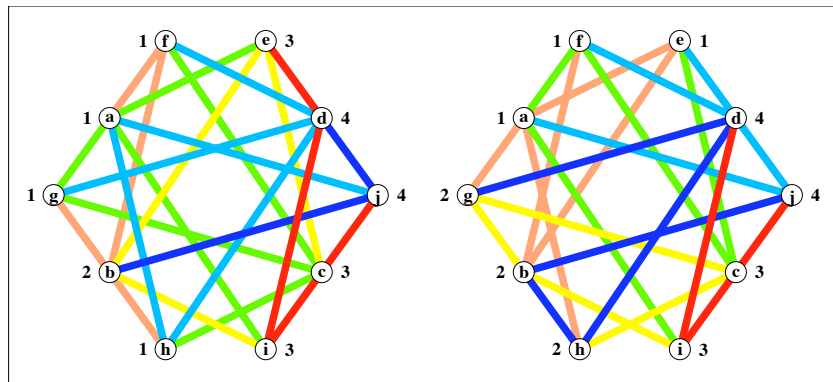


Figure: Vertex packets are trees (l), paths (r).

# Eliminate Obstacles - Eliminate Shelling

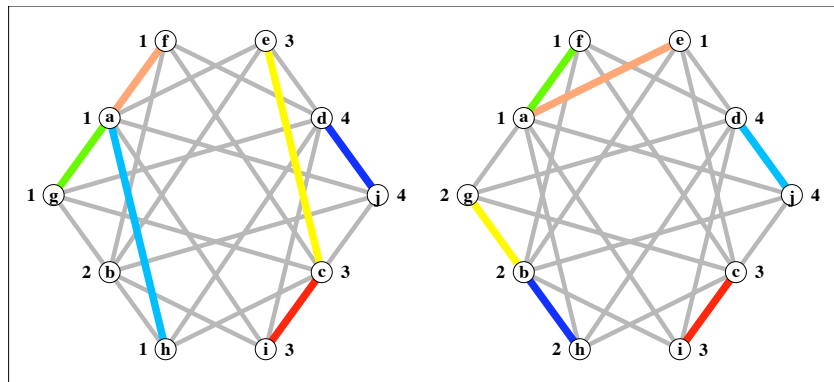


Figure: Vertex packets are trees (l), paths (r).

## Understanding these drawings:

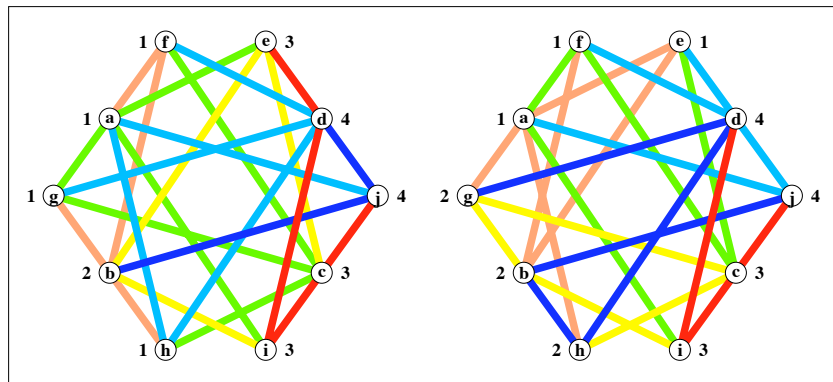


Figure: These drawings may seem complicated, but are easily analyzed.

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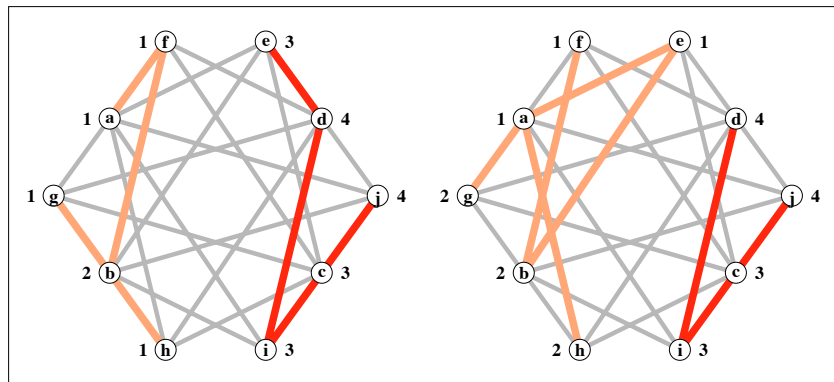


Figure: Trees  $T_{12}$ ,  $T_{34}$ .

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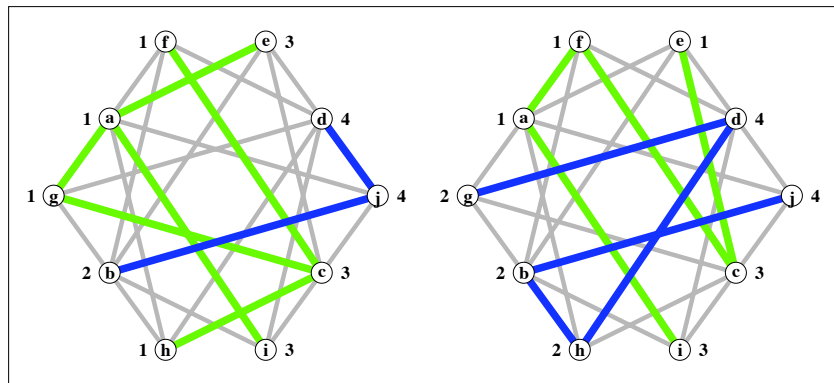


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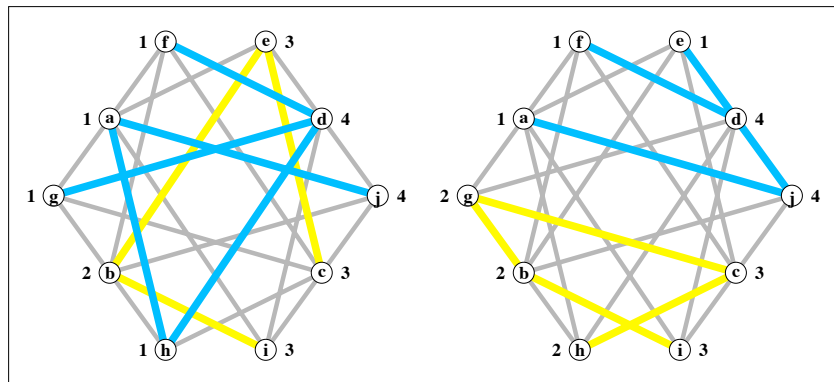


Figure: Trees  $T_{14}$ ,  $T_{23}$ .

The vertex set can not always be partitioned into paths.

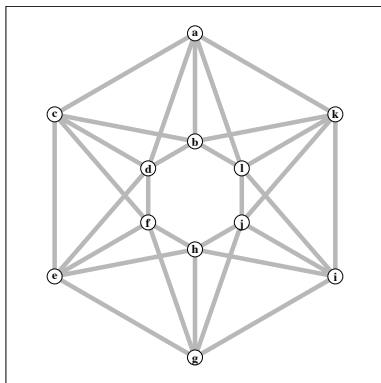
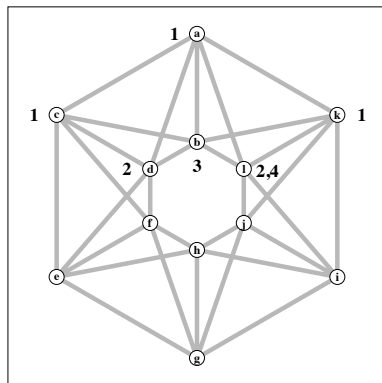


Figure: A hinged ring of tetrahedra.

3-isostatic graphs do not necessarily have  
maps to  $K_4$  in which

( $\mathcal{P}$ ) induced graphs on vertex packets are paths.

The vertex set can not always be partitioned into paths.



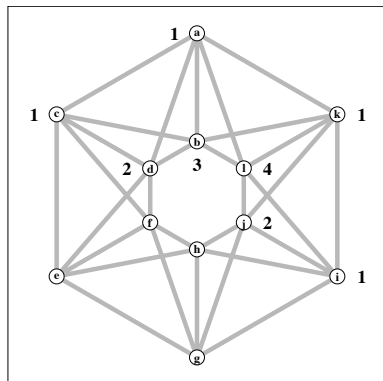
There must be a path of length  $\geq 3$ ,  
not within a single tetrahedron.

The vertex  $b$  is isolated with its image 3.

There must be a path of length  $\geq 4$ .



The vertex set can not always be partitioned into paths.



$l$  must be 4, otherwise there is no 2-path from  $d$  to  $l$ .  
Then values 3 and 4 are isolated at  $b$  and  $l$ ,  
So only 1 and 2 are available for tetrahedron  $efgh$ .

# Freely shellable semi-simplicial maps

In practice, freely-shellable maps seem to abound,  
and seem much easier to find “by hand”  
than more general maps  
for which you must check shellability.

# Freely shellable semi-simplicial maps

What is more, freely-shellable maps  
have relatively few loops that need to be assigned.

# Partitions that Produce Freely Shellable Maps

To prove a graph  $G(V, E)$  is isostatic, it suffices to exhibit a partition  $\pi$  of the vertex set  $V$  having three properties  $\mathcal{P}_i$  (see below).

The main criterion  $\mathcal{P}_3$  is Richard Rado's matroid basis matching condition.

# Partitions that Produce Freely Shellable Maps

## Theorem: Rado's Basis Matching Theorem

Given any relation  $R$  from a set  $X$  to a set  $S$  of elements of a matroid  $M(S)$ , then there is matching in  $R$  from  $X$  to a basis for the matroid  $M(S)$  if and only if

- the cardinality  $|X| = \text{rank } \rho(S)$  of the matroid  $M$ ,
- and, for every subset  $A \subset X$ ,
- the cardinality  $|A| \leq \rho(A)$ ,
- the rank of its image  $R(A)$  in  $M(S)$ .

# Bibliography on Matroid Matching

Richard Rado,

*A Theorem on Independence Relations,*

Quarterly J. of mathematics, Oxford **13** (1942), 83-89.

# Bibliography on Matroid Matching

Joseph P. S. Kung, Gian-Carlo Rota, Catherine H. Yan,  
*Combinatorics: The Rota Way*,  
Cambridge University Press, 2009.

# Bibliography on Matroid Matching

Kazuo Murota,  
*Matrices and Matroids for Systems Analysis*  
Springer Verlag,  
*Algorithms and Combinatorics***20** (2000),(revised 2010).



## Bibliography on Matroid Matching

And an article which led us to the possibility of insisting that vertex packets induce paths:

Roger K. S. Poh,

*On the Linear Vertex-Arboricity of a Planar Graph*

Journal of Graph Theory, **14 No. 1** (1990), 73-75.

# A Matroid Union

Given a partition of the vertex set of  $G$ ,  
define bridges and loops,  
and for each  $ij$  construct the matroid minor:  
restrict to the induced subgraph  
on the union of the two packets,  
and contract by its bridges.

Then take the matroid union over all pairs  $ij$

# A Matroid Union

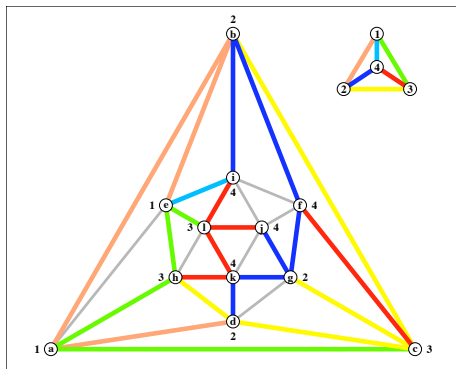


Figure: The bridges of a map on the icosahedron.

# A Matroid Union

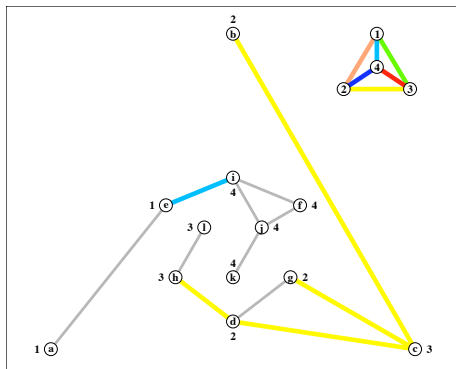


Figure: Restrictions to packet unions  $V_1 \cup V_4$  and  $V_2 \cup V_3$ .

# Characterization of Partitions for Freely-Shellable Maps

**Theorem:** A partition  $\pi$  of the vertex set of a graph  $G(V, E)$  is the inverse image partition of a freely-shellable semi-simplicial map  $f : G \rightarrow K_{d+1}$  if and only if the partition  $\pi$  has the following three properties  $\mathcal{P}_i$

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( $\mathcal{P}_1$ ) The induced subgraph  $G_i$  on any part  $\pi_i$  of  $\pi$  is independent (circuit-free).

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- $(\mathcal{P}_2)$  For any pair  $ij$ , the bridge subgraph  $G(V_i \cup V_j, B_{ij})$  is independent.
- $(\mathcal{P}_3)$  The relation  $\mathcal{R}$  between the set of loops of  $G$  and the set of elements of the matroid union  $M$  satisfies the Rado condition for basis matching:  
 $|L| = \rho(M)$  and

$$\forall A \subseteq E, |A| \leq \rho(\mathcal{R}(A)).$$



# A Partition Not Satisfying the Rado Condition

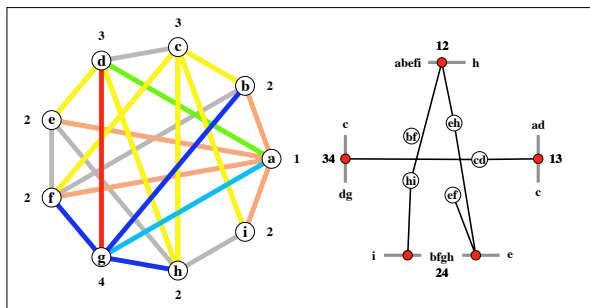


Figure: Partition  $(a)(befhi)(cd)(g)$  does not satisfy the Rado condition.

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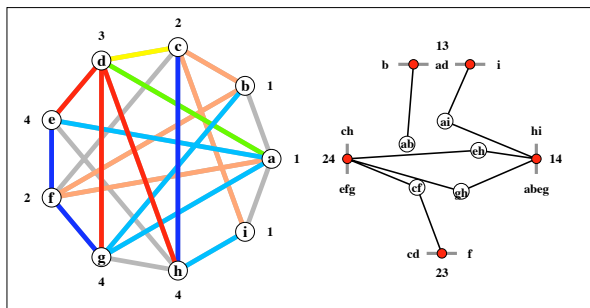


Figure: Partition  $(a)(befhi)(cd)(g)$  has 2 compatible loop maps.

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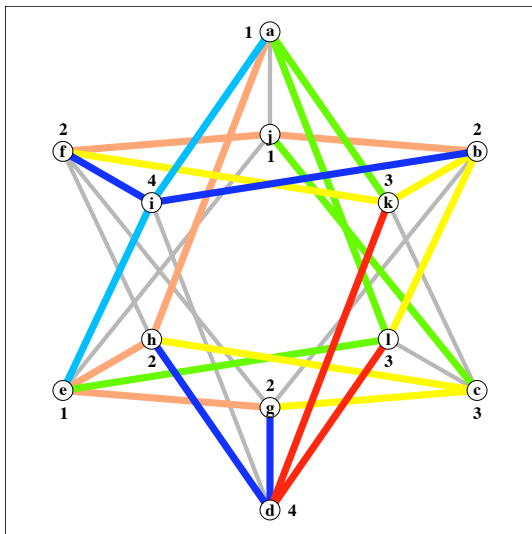


Figure: A non-Rado partition for  $K_{6,6}$  less 6 edges. (edge  $di$ !)



# A Partition Satisfying the Rado Condition

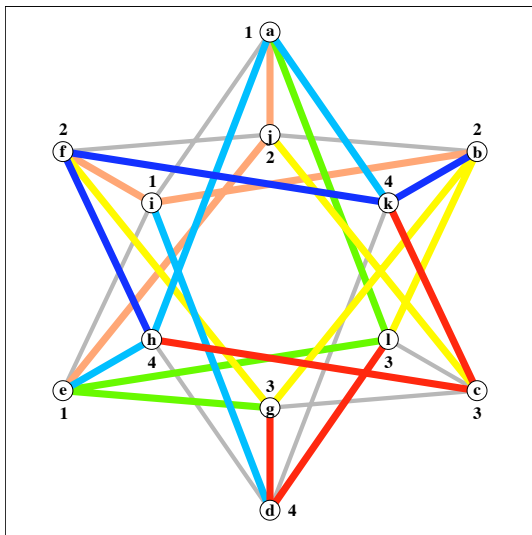


Figure: A partition with 32 compatible loop maps.

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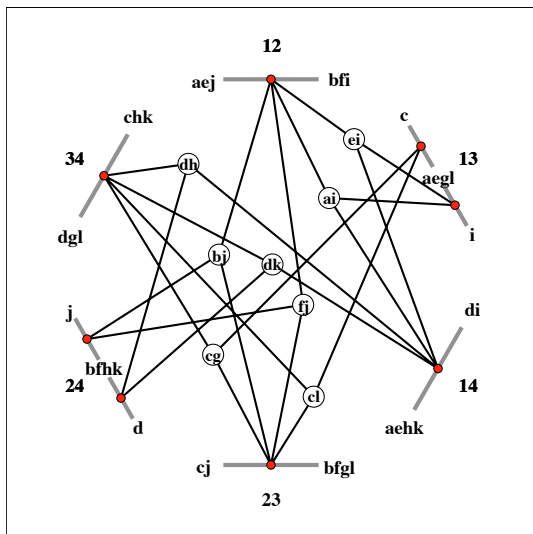


Figure: The Rado relation  $\mathcal{R}$  for that partition.

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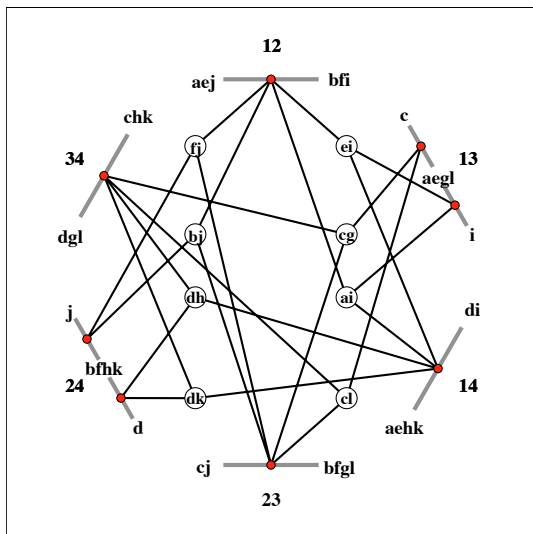


Figure: The symmetry of  $\mathcal{R}$  is perhaps more visible here.





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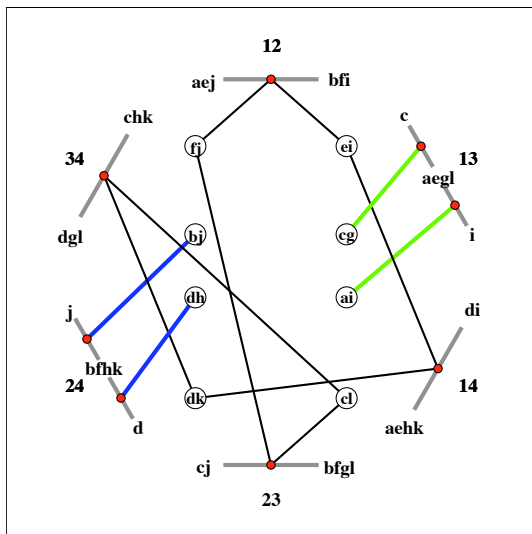


Figure: After four independent binary choices, a cycle remains.

# The Road Ahead

It remains to prove that any 3-isostatic graph has a freely-shellable semi-simplicial map to the simplex  $K_4$ .

# The Road Ahead

This has always been the hard part of the problem!

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Followed by a rapid retreat  
from an untenable position!

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Which properties of isostatic graphs  
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Which properties of isostatic graphs might permit us to prove the conjecture?

We lean toward an analogue in  $d = 3$   
of Tay's proof for  $d = 2$ .

## Toward an Analogue of Tay's Proof for $d = 2$

We use the  $(3v - 6) \times 6v$   
projective rigidity matrix  $R$ ,  
and the  $(3v + 6) \times 6v$  matrix  $S$   
whose rows span the orthogonal complementary subspace.

## Toward an Analogue of Tay's Proof for $d = 2$

By **Hodge star** complementation,  
the determinants of full-size minors of  $R$   
are equal to the determinants  
of the complementary full-size minors of  $S$   
up to a sign  $\pm 1$  of the bipartition of the column set,  
and up to a fixed polynomial quantity  $Q$ ,  
called the **pure condition** or **resolving bracket**,  
which is non-zero exactly when the graph is isostatic.

## Toward an Analogue of Tay's Proof for $d = 2$

The column matroids of  $R$  and of  $S$  are dual to one another,  
and are independent of the graph  $G$  in question!

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The column matroids of  $R$  and of  $S$  are dual to one another,  
and are independent of the graph  $G$  in question!

( $Q \neq 0$  exactly when the rows of  $R$  form a basis for  
the space of external equilibrium loads  
on the set  $V$  of vertices of  $G$ ,  
regarded as a single rigid body.)

# The Orthogonal Complementary Matrices $S$ and $R$

a12	b12	c12	d12	e12	a13	b13	c13	d13	e13	a14	b14	c14	d14	e14	a23	b23	c23	d23	e23	a24	b24	c24	d24	e24	a34	b34	c34	d34	e34			
1	1	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	C12						
0	0	0	0	0	1	1	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	C13						
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	0	0	0	0	0	0	0	0	0	0	0	C14						
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	0	0	0	0	0	0	C23						
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	0	C24						
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	C34						
a3	0	0	0	0	-a2	0	0	0	0	0	0	0	0	0	a1	0	0	0	0	0	0	0	0	0	0	a123						
a4	0	0	0	0	0	0	0	0	0	-a2	0	0	0	0	0	0	0	0	0	a1	0	0	0	0	0	0	a124					
0	0	0	0	0	a4	0	0	0	0	-a3	0	0	0	0	0	0	0	0	0	0	0	0	0	a1	0	0	a134					
0	b3	0	0	0	0	-b2	0	0	0	0	0	0	0	0	b1	0	0	0	0	0	0	0	0	0	0	0	b123					
0	b4	0	0	0	0	0	0	0	0	0	-b2	0	0	0	0	0	0	0	0	b1	0	0	0	0	0	0	b124					
0	0	0	0	0	0	b4	0	0	0	0	-b3	0	0	0	0	0	0	0	0	0	0	0	0	b1	0	0	b134					
0	0	c3	0	0	0	-c2	0	0	0	0	0	0	0	0	c1	0	0	0	0	0	0	0	0	0	0	0	c123					
0	0	c4	0	0	0	0	0	0	0	0	-c2	0	0	0	0	0	0	0	0	c1	0	0	0	0	0	0	c124					
0	0	0	0	0	0	0	c4	0	0	0	0	-c3	0	0	0	0	0	0	0	0	0	0	0	c1	0	0	c134					
0	0	0	d3	0	0	0	0	0	0	0	0	0	0	0	0	0	d1	0	0	0	0	0	0	0	0	0	d123					
0	0	0	d4	0	0	0	0	0	0	0	0	0	-d2	0	0	0	0	0	0	0	0	0	d1	0	0	0	d124					
0	0	0	0	0	0	0	0	0	0	0	0	0	d4	0	0	0	0	0	0	0	0	0	0	0	d1	0	d134					
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	d123				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	d124				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	d134				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	e123				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	e124				
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	e134				
ac12	0	-ac12	0	0	ac13	0	-ac13	0	0	ac14	0	-ac14	0	0	ac23	0	-ac23	0	0	ac24	0	-ac24	0	0	ac34	0	-ac34	0	0	ac		
ad12	0	0	-ad12	0	ad13	0	0	-ad13	0	ad14	0	0	-ad14	0	ad23	0	0	-ad23	0	ad24	0	0	-ad24	0	0	ad34	0	0	-ad34	0	ad	
ae12	0	0	0	0	-ae12	0	0	0	0	-ae13	0	0	0	0	-ae23	0	0	0	0	-ae24	0	0	0	0	-ae34	0	0	0	0	-ae		
0	bc12	-bc12	0	0	0	bc13	-bc13	0	0	0	bc14	-bc14	0	0	0	bc23	-bc23	0	0	0	bc24	-bc24	0	0	0	bc34	-bc34	0	0	bc		
0	bd12	0	-bd12	0	0	bd13	0	-bd13	0	0	bd14	0	-bd14	0	0	bd23	0	-bd23	0	0	bd24	0	-bd24	0	0	bd34	0	-bd34	0	0	bd	
0	bc12	0	0	-bc12	0	bc13	0	0	-bc13	0	bc14	0	0	-bc14	0	bc23	0	0	-bc23	0	bc24	0	0	-bc24	0	bc34	0	0	-bc34	0	bc	
0	0	cd12	-cd12	0	0	cd13	-cd13	0	0	0	cd14	-cd14	0	0	0	cd23	-cd23	0	0	0	cd24	-cd24	0	0	0	cd34	-cd34	0	0	cd		
0	0	ce12	0	0	-ce12	0	0	0	0	0	ce13	0	0	0	0	ce23	0	0	0	0	ce24	0	0	0	0	ce34	0	0	0	-ce34	ce	
0	0	0	de12	-de12	0	0	0	0	0	0	0	de13	-de13	0	0	0	0	0	0	0	0	de24	-de24	0	0	0	0	de34	-de34	0	0	de

Figure: Columns grouped by trees  $T_{ij}$ .

# The Orthogonal Complementary Matrices $S$ and $R$

a12 a13 a14 a23 a24 a34	b12 b13 b14 b23 b24 b34	c12 c13 c14 c23 c24 c34	d12 d13 d14 d23 d24 d34	e12 e13 e14 e23 e24 e34	
1 0 0 0 0 0	1 0 0 0 0 0	1 0 0 0 0 0	1 0 0 0 0 0	1 0 0 0 0 0	C12
0 1 0 0 0 0	0 1 0 0 0 0	0 1 0 0 0 0	0 1 0 0 0 0	0 1 0 0 0 0	C13
0 0 1 0 0 0	0 0 1 0 0 0	0 0 1 0 0 0	0 0 1 0 0 0	0 0 1 0 0 0	C14
0 0 0 1 0 0	0 0 0 1 0 0	0 0 0 1 0 0	0 0 0 1 0 0	0 0 0 1 0 0	C23
0 0 0 0 1 0	0 0 0 0 1 0	0 0 0 0 1 0	0 0 0 0 1 0	0 0 0 0 1 0	C24
0 0 0 0 0 1	0 0 0 0 0 1	0 0 0 0 0 1	0 0 0 0 0 1	0 0 0 0 0 1	C34
a3 -a2 0 a1 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	a123
a4 0 -a2 0 a1 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	a124
0 -a4 -a3 0 0 a1	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	a134
0 0 0 0 0 0	b3 -b2 0 b1 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	b123
0 0 0 0 0 0	b4 0 -b2 0 b1 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	b124
0 0 0 0 0 0	0 b4 -b3 0 0 b1	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	b134
0 0 0 0 0 0	0 0 0 0 0 0	c3 -c2 0 c1 0 0	0 0 0 0 0 0	0 0 0 0 0 0	c123
0 0 0 0 0 0	0 0 0 0 0 0	c4 0 -c2 0 c1 0	0 0 0 0 0 0	0 0 0 0 0 0	c124
0 0 0 0 0 0	0 0 0 0 0 0	0 c4 -c3 0 0 c1	0 0 0 0 0 0	0 0 0 0 0 0	c134
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	d3 -d2 0 d1 0 0	0 0 0 0 0 0	d123
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	d4 0 -d2 0 d1 0	0 0 0 0 0 0	d124
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 d4 -d3 0 0 d1	0 0 0 0 0 0	d134
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	e3 -c2 0 c1 0 0	e123
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	e4 0 -c2 0 c1 0	e124
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	0 e4 -c3 0 0 c1	e134
ae12 ae13 ae14 ae23 ae24 ae34	0 0 0 0 0 0	-ae12 -ae13 -ae14 -ae23 -ae24 -ae34	0 0 0 0 0 0	0 0 0 0 0 0	ae
ad12 ad13 ad14 ad23 ad24 ad34	0 0 0 0 0 0	0 0 0 0 0 0	-ad12 -ad13 -ad14 -ad23 -ad24 -ad34	0 0 0 0 0 0	ad
ae12 ae13 ae14 ae23 ae24 ae34	0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	-ae12 -ae13 -ae14 -ae23 -ae24 -ae34	ae
0 0 0 0 0 0	be12 be13 be14 be23 be24 be34	-be12 -be13 -be14 -be23 -be24 -be34	0 0 0 0 0 0	0 0 0 0 0 0	be
0 0 0 0 0 0	bd12 bd13 bd14 bd23 bd24 bd34	0 0 0 0 0 0	-bd12 -bd13 -bd14 -bd23 -bd24 -bd34	0 0 0 0 0 0	bd
0 0 0 0 0 0	be12 be13 be14 be23 be24 be34	0 0 0 0 0 0	0 0 0 0 0 0	-be12 -be13 -be14 -be23 -be24 -be34	be
0 0 0 0 0 0	0 0 0 0 0 0	cd12 cd13 cd14 cd23 ce24 cd34	-cd12 -cd13 -cd14 -cd23 -cd24 -cd34	0 0 0 0 0 0	cd
0 0 0 0 0 0	0 0 0 0 0 0	ce12 ce13 ce14 ce23 ce24 ce34	0 0 0 0 0 0	-ce12 -ce13 -ce14 -ce23 -ce24 -ce34	ce
0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0	de12 de13 de14 de23 de24 de34	-de12 -de13 -de14 -de23 -de24 -de34	de

Figure: Columns grouped by vertices  $v$ .

# The Orthogonal Complementary Matrices $S$ and $R$

Any set of columns in  $R$  labeled by a single vertex,  
say by  $a$   
and by a circuit in  $K_4$ ,  
such as 12, 23, 34, 14,  
are dependent.



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are dependent.

Any set of columns in  $R$  labeled by a edge of  $K_4$ ,  
say by 12  
and by all vertices  $a$ ,  
are dependent.

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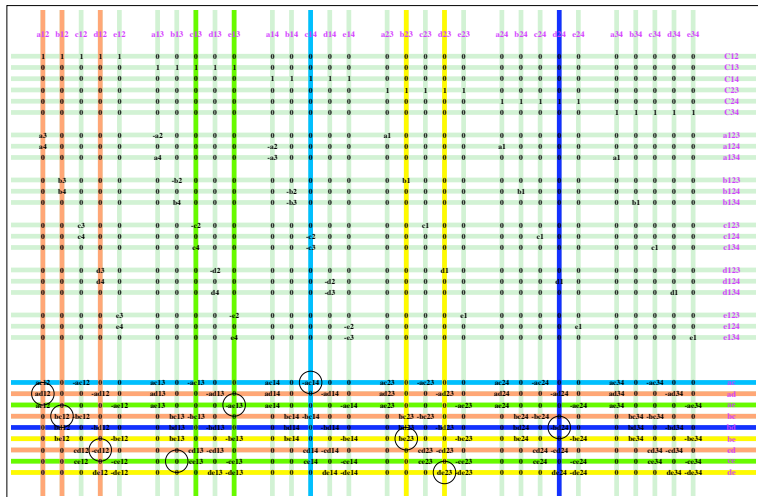


Figure: From a non-zero diagonal to a (rooted) freely shellable map.

# The Orthogonal Complementary Matrices $S$ and $R$

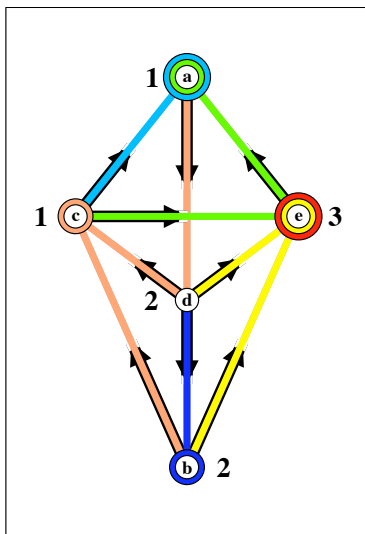


Figure: The corresponding rooting of a freely shellable map.

# An analogue of Henneberg reduction?

Is it possible to reduce any isostatic graph to an isostatic graph on one fewer vertex, by a procedure that, when repeated, leads, step-by-step, to a map?

Grazie

Thank you for your attention.

This paper should be up on the [arXiv](#) soon:

*Isostatic Structures:  
Using Richard Rado's Independent Matchings*