

# Partial Word Representation

F. Blanchet-Sadri

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# Partial words and compatibility

- ▶ A partial word is a sequence that may have undefined positions, called holes and denoted by  $\diamond$ 's, that match any letter of the alphabet  $A$  over which it is defined (a full word is a partial word without holes); we also say that  $\diamond$  is compatible with each  $a \in A$ .

$a\diamond b\diamond aab$  is a partial word with two holes over  $\{a, b\}$

- ▶ Two partial words  $w$  and  $w'$  of equal length are compatible, denoted by  $w \uparrow w'$ , if  $w[i] = w'[i]$  whenever  $w[i], w'[i] \in A$ .

	$a$	$\diamond$	$b$	$\diamond$	$a$		$a$	$\diamond$	$b$	$\diamond$	$a$
$\uparrow$											
	$\diamond$	$\diamond$	$b$	$a$	$a$		$\diamond$	$\diamond$	$a$	$a$	$a$

# Factor and subword

- ▶ A partial word  $u$  is a **factor** of the partial word  $w$  if  $u$  is a block of consecutive symbols of  $w$ .

$\diamond a \diamond$  is a factor of  $aa \diamond a \diamond b$

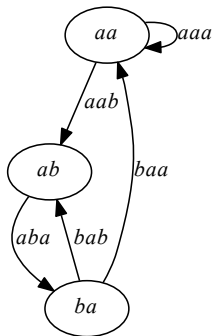
- ▶ A full word  $u$  is a **subword** of the partial word  $w$  if it is compatible with a factor of  $w$ .

$aaa, aab, baa, bab$  are the subwords of  $aa \diamond a \diamond b$   
corresponding to the factor  $\diamond a \diamond$

## Some computational problems

- ▶ We define **REP**, or the problem of deciding whether a set  $S$  of words of length  $n$  is **representable**, i.e., whether  $S = \text{sub}_w(n)$  for some integer  $n$  and partial word  $w$ .
- ▶ If  $h$  is a non-negative integer, we also define  **$h$ -REP**, or the problem of deciding whether  $S$  is  **$h$ -representable**, i.e., whether  $S = \text{sub}_w(n)$  for some integer  $n$  and partial word  $w$  with exactly  $h$  holes.

Rauzy graph of  $S = \{aaa, aab, aba, baa, bab\}$



$S$  is 0-representable by  $w = aaababaa$

# Why partial words? (Compression of representations)

$$S = \{aaa, aab, aba, baa, bab\}$$

is representable by

*aaababaa*

# Why partial words? (Compression of representations)

$$S = \{aaa, aab, aba, baa, bab\}$$

is representable by

◇ *aabab*

# Why partial words? (Compression of representations)

$$S = \{aaa, aab, aba, baa, bab\}$$

is representable by

$a \diamond a \diamond$



# Why partial words? (Representation of non-0-representable sets)

- ▶ Set  $S$  of 30 words of length six:

1	aaaaaa	6	aabbba	11	abbbba	16	baabbb	21	bbabab	26	bbbabb
2	aaaaab	7	aabbba	12	abbbab	17	bababb	22	bbabbb	27	bbbbaa
3	aaaabb	8	aabbbb	13	abbbba	18	babbba	23	bbbbaa	28	bbbbaa
4	aaabba	9	ababbb	14	abbbbb	19	babbbb	24	bbbaab	29	bbbbaa
5	aaabbb	10	abbaab	15	baabba	20	bbaabb	25	bbbaba	30	bbbbaa

- ▶ Rauzy graph  $(V, E)$  of  $S$ , where  $E = S$  and  $V = \text{sub}_S(5)$  consists of 20 words of length five:

1	aaaaa	5	aabbb	9	abbbb	13	bbaaa	17	bbbba
2	aaaab	6	ababb	10	baabb	14	bbbaa	18	bbbba
3	aaabb	7	abbaa	11	babab	15	bbaba	19	bbbba
4	aabba	8	abbbba	12	babbb	16	bbabb	20	bbbba





# Membership of $h$ -REP in $\mathcal{P}$

## Theorem

$h$ -REP is in  $\mathcal{P}$  for any fixed non-negative integer  $h$ .

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F. Blanchet-Sadri and S. Simmons, Deciding representability of sets of words of equal length. *Theoretical Computer Science* **475** (2013) 34–46

# Membership of REP in $\mathcal{P}$

## Theorem

REP is in  $\mathcal{P}$ .

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F. Blanchet-Sadri and S. Munteanu, Deciding representability of sets of words of equal length in polynomial time. Submitted

# Open problems

- ▶ Characterize the sets of words that are representable.
- ▶ Characterize minimal representing partial words (can they be constructed efficiently?)