## A NEW COMPLEXITY MEASURE FOR WORDS BASED ON PERIODICITY

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## Periods of a word

 $w = a_1 a_2 \dots a_n$ 

A positive integer  $p \le |w|$  is a period of w if

$$a_{i+p} = a_i$$
, for  $i = 1, 2, ..., n-p$ 

The smallest period of w is called the period of w and is denoted by p(w)

abaababaaba has periods 5 and 8

## Local periods

 $w = a_1 a_2 \dots a_n$ 

A non-empty word u is a repetition of w at the point i if w = xy, with |x| = i and the following holds:

$$\mathsf{A}^* \, x \ \cap \ \mathsf{A}^* u \ \neq \varphi \quad \text{and} \quad y \ \mathsf{A}^* \ \cap \ u \ \mathsf{A}^* \neq \varphi$$

The local period of w at the point i is:

p(w,i) = min { |u| : u is a repetition of w at the point i }

An example of repetition and local period w = a b a a b a a b a a b 1 2 3 4 5 6 7 8 9 10 11 12 8 1 p(w,3) = 1p(w,7) = 8

a b a a b a b a a b a a b 2 3 1 5 2 2 8 1 3 3 1 3

A point i is critical if p(w,i) = p(w)

Critical Factorization Theorem (CFT) (Cesari-Vincent, 1978; Duval, 1979) If  $|w| \ge 2$ , in any sequence of m = max {1, p(w)-1} consecutive points there is a critical one, i.e. there exists a positive integer i such that p(w,i) = p(w).

A point i is called left external if i < p(w,i). From CFT, the first critical point is left external.

## Local periods in infinite words

Theorem. An infinite words is recurrent if and only if at any point there is a repetition

**Periodicity function** of an infinite recurrent word x:

 $p_x(n) = min \{ |u| : u \text{ is a repetition at the point } n \}$ 

Theorem. An infinite recurrent word x is periodic if and only if the periodicity function  $p_x$  is bounded. Moreover  $p(x) = \sup \{p_x(n) : n \ge 1\}$ 

## Gap Theorem

Theorem. Let x be an infinite recurrent word. Then either  $p_x$  is bounded, i.e. x is periodic, or  $p_x(n) \ge n+1$ , for infinitely many integers n.

Analogous to the Coven-Hedlund theorem:

Theorem (Coven-Hedlund). The (factor) complexity function  $c_x$  of an infinite word x either is bounded, and in such a case x is periodic, or  $c_x(n) \ge n+1$ , for all integers n





Thue-Morse

#### Fibonacci



# Characteristic Sturmian words are extremal for the CFT

Theorem. Let x be an infinite recurrent word. X is a characteristic sturmian word if and only if  $p_x(n) \le n + 1$  for all  $n \ge 1$  and  $p_x(n) = n + 1$  for infinitely many integers n.

Equivalently:

The characteristic sturmian words are exactly the recurrent non periodic words x such that  $p_x(n) \le n + 1$ .

## Finite Standard words

Let  $q_0$ ,  $q_1$ ,  $q_2$ , ..... be a sequence of non-negative integers , with  $q_i > 0$  for i > 0.

Consider the sequence of words  $\{s_n\}_{n\geq 0}$  defined as follows:

$$s_0 = b$$
  
 $s_1 = a$   
 $s_{n+1} = s_n^{q_{n-1}} s_{n-1}$ 

## Characteristic Sturmian words

The sequence  $\{s_n\}_{n\geq 0}$  converges to a limit x that is an infinite characteristic Sturmian word.

The sequence  $\{s_n\}_{n\geq 0}$  is called the approximating sequence of x and  $(q_0, q_1, q_2, ...)$  is the directive sequence of x.

Each finite word  $s_n$  is called a standard word and it is univocally determined by the (finite) directive sequence  $(q_0, q_1, ..., q_{n-2})$ . Computation of the periodicity function of a characteristic Sturmian word

If **x** is (the Fibonacci) a characteristic Sturmian word, then the function  $p_x(n)$  can be computed from the (Zeckendorf) Ostrowski representation of the integer n+1

(J. Shallit, L. Schaeffer)

## Non-characteristic Sturmian words

Remark that the characterization theorem holds true just for characteristic Sturmian words, not for all Sturmian words:

y = aababaabaabaabaab....  
$$p_v(2) = 5$$
  $p_v(5) = 8$ 

#### Theorem

The periodicity function characterizes any finite or infinite binary word up to exchange of letters.

Remark: this is not true in alphabeths having more than two letters.

b c a c b c a b b 1 8 8 8 8 8 8 8 1 b b c a c b a c b b

## **Periodicity Complexity**

The periodicity function has a strong fluctuation, and this is not convenient for certain purposes.

So, we introduce the periodicity complexity function  $h_x(n)$  of an infinite word x, defined as follows:

$$h_x(n) = \frac{1}{n} \sum_{j=1}^n p_x(j)$$

Theorem

If x is an infinite periodic word, then the periodicity complexity function  $h_x(n)$  is bounded.

The converse is not true:

There exist non-periodic recurrent words having bounded periodicity complexity.

A non-periodic word with bounded periodicity complexity

Consider a sequence of finite words recursively defined as follows:

- $w_0 = ab$  $w_{n+1} = w_n a^{2|w_n|} w_n$
- $w_1 = abaaaaab$

 $\mathbf{w} = \lim_{n \to \infty} w_n$ 

**Theorem** lim sup  $h_w(n) = \sup h_w(n) = 7$ 

## The Fibonacci word

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## Theorem $h_f(n)$ grows as $\Theta(\log n)$

## The Thue-Morse word

#### **t** = abbabaabbaababbaabbaabba......

#### Theorem $h_t(n)$ grows as $\Theta(n)$

An infinite recurrent word with arbitrary high periodicity complexity

Let  $v_n$  be the finite binary word obtained by concatenating in the lexicographic order all the words of length n.

 $v_1 = ab$ 

- $v_2 = aaabbabb$
- $v_3 = aaaaababaabbbaababbbabbb$

For any function **f** from N to N consider the sequence of words:  $z_1 = v_1$  $z_{n+1} = z_n b z_n^{[2 f(|z_n|+1)]} v_{n+1}$ 

#### Consider the infinite word $\mathbf{z} = \lim z_n$

**Theorem** For infinitely many j,  $h_z(j) > f(j)$ .

# Problems

 Does there exist a uniformly recurrent nonperiodic word having bounded periodicity complexity ?

• Does there exist a uniformly recurrent word with arbitrary high periodicity complexity ?

• Evaluate the periodicity complexity of other special words